# Query Optimization 

Exercise Session 11

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July 7, 2014

## Homework

- Last Homework: Quick Pick
- Union-Find
- Union-By-Size
- Path-Shortening


## Quick-Pick this!



| Edge | Roots | Union-Find Array |
| :---: | :---: | :---: |
| $(D, C)$ |  |  |
| $(A, B)$ |  |  |
| $(E, F)$ |  |  |
| $(E, D)$ |  |  |
| $(A, D)$ |  |  | Join Trees

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| :---: | :--- | :--- |
| $(D, C)$ | $D, C$ | $\left[-\frac{B}{-},-\frac{C}{-}, 2,-\frac{E}{-}\right]$ |
| $(A, B)$ |  |  |
| $(E, F)$ |  |  |
| $(E, D)$ |  |  |
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| $(A, B)$ |  |  |  |
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| :---: | :--- | :---: | :---: |
| $(D, C)$ | $D, C$ | $\left[-\frac{A}{-}, \frac{C}{-}, \frac{D}{2}, \frac{E}{-},-\right]$ | $\{A, B, C \bowtie D, E, F\}$ |
| $(A, B)$ | $A, B$ | $\left[-\frac{A}{-}, \frac{C}{-}, \frac{E}{-},-\frac{-}{-}\right]$ |  |
| $(E, F)$ |  |  |  |
| $(E, D)$ |  |  |  |
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| $(D, C)$ | $D, C$ | $\left[-\frac{B}{-}, \frac{C}{D}, \frac{E}{-}, \frac{F}{-}\right]$ | $\{A, B, C \bowtie D, E, F\}$ |
| $(A, B)$ | $A, B$ | $\left[-\frac{A}{-}, \frac{D}{-}, \frac{E}{-}, \frac{-}{-}\right]$ | $\{A \bowtie B, C \bowtie D, E, F\}$ |
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| $(A, B)$ | $A, B$ | $\left[-\frac{A}{-}, \frac{D}{-}, \frac{E}{F}, \frac{-}{-}\right]$ | $\{A \bowtie B, C \bowtie D, E, F\}$ |
| $(E, F)$ | $E, F$ | $\left[-A, 0,-\frac{B}{-}, \frac{E}{-}, 4\right]$ | $\{A \bowtie B, C \bowtie D, E \bowtie F\}$ |
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| $(E, D)$ | $E, C$ |  | $\{A \bowtie B,(C \bowtie D) \bowtie(E \bowtie F)\}$ |
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| $(E, F)$ | $E, F$ |  | $\{A \bowtie B, C \bowtie D, E \bowtie F\}$ |
| $(E, D)$ | E, C | $\left[\begin{array}{ccc} A & B & C \\ - & D & D \\ 0, & E & E \\ 2, & 2, & 4 \end{array}\right]$ | $\{A \bowtie B,(C \bowtie D) \bowtie(E \bowtie F)\}$ |
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| $(A, D)$ | $A, C$ | $\left[\begin{array}{lll} A \\ - & B, & C \\ 0 & D, & D, \\ \hline \end{array}, 4\right]$ | $\{(A \bowtie B) \bowtie((C \bowtie D) \bowtie(E \bowtie F))\}$ |

A: A branch of mathematics concerning the study of finite or countable discrete structures.
Q: What is ?

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Q: What is combinatorics?

## Combinatorics 101

Given a set of $n$ elements, how many distinct $k$-element subsets can be formed?

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$$
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$$

Example: Choose 3 out of $5:\binom{5}{3}=\frac{5!}{2!\cdot 3!}=\frac{120}{2 \cdot 6}=10$


## Direct, Uniform, Distinct

## Waters/Yao Bottom-Up

Given $m$ pages with $n$ tuples on each page, e.g. a total of $N=m \cdot n$ tuples:


- How many distinct subsets of size $k$ exist?


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- How many distinct subsets of size $k$ exist, where a page does not contain any chosen tuples? Choose $k$ from all but one page, i.e. from $N-n$ tuples:


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So the probability that a page contains none of the $k$ tuples is


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So the probability that a page contains none of the $k$ tuples is

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p:=\frac{\binom{N-n}{k}}{\binom{N}{k}}
$$

- What is the probability that a certains page contains at least one tuple?


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$$

- What is the probability that a certains page contains at least one tuple? $1-p \ldots$ unless all pages have to be involved ( $k>N-n$ ).
- Multiplied by the number of pages, we get the number of qualifying pages, denoted $\overline{\mathcal{Y}}_{n}^{N, m}(k)$.


## Approximation

Let $m=50, n=1000 \Rightarrow N=50 k, k=100$
Yao (exact) : $p=\frac{\binom{N-n}{k}}{\binom{N}{k}}=\prod_{i=0}^{k-1} \frac{N-n-i}{N-i}=\prod_{i=0}^{99} \frac{49 k-i}{50 k-i}=13.2 \%$ Waters : $p \approx\left(1-\frac{k}{N}\right)^{n}$

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Waters : $p \approx\left(1-\frac{k}{N}\right)^{n} \approx 13.5 \%$

## Direct, Uniform, Non-Distinct

## Combinatorics 101 revisited

- Now with replacement: How many distinct multisets exist chosing $k$ from $n$ ?


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## Combinatorics 101 revisited

- Now with replacement: How many distinct multisets exist chosing $k$ from $n$ ? As many as there are distinct sets chosing $k$ from $n+k-1$ !
- Bijection between multisets and sets. From multiset to set: $f:\left(x_{1}, x_{2}, \ldots, x_{k}\right) \mapsto\left(x_{1}+0, x_{2}+1, \ldots, x_{k}+(k-1)\right)$


## Combinatorics 101 revisited

- Now with replacement: How many distinct multisets exist chosing $k$ from $n$ ? As many as there are distinct sets chosing $k$ from $n+k-1$ !
- Bijection between multisets and sets. From multiset to set: $f:\left(x_{1}, x_{2}, \ldots, x_{k}\right) \mapsto\left(x_{1}+0, x_{2}+1, \ldots, x_{k}+(k-1)\right)$
- Example: Choose 2 from 4
- \# sets: $\binom{4}{2}$
- \# multisets: $\binom{4+2-1}{2}$



## Cheung

- Like Yao, but not necessarily distinct
- Same formula as Yao, but:
- We don't need to distinguish cases when computing the probability that a bucket contains at least one item
- We substitue $N$ by $N+k-1$ to compute $\tilde{p}$

Direct, Non-Uniform, Distinct

## Direct, Non-Uniform, Distinct



Assume that $n_{j}>0 \forall j \in[1, m]$, then the expected number of qualifying pages is

$$
\sum_{j=1}^{m}\left(1-\frac{\binom{N-n_{j}}{k}}{\binom{N}{k}}\right)
$$

With $N=\sum_{j=1}^{m} n_{j}$.

## Distribution Function

- The number of possibilities to select $x\left(x \leq n_{j}\right)$ items from bucket $j$ is $\binom{n_{j}}{x}$.
- The number of possibilities to draw the remaining $k-x$ items from other buckets is $\binom{N-n_{j}}{k-x}$.
- Recall: The number of possibilities to draw $k$ items from $N$ is $\binom{N}{k}$.
$\Rightarrow$ The probability that $x$ items qualify from bucket $j$ is

$$
\frac{\binom{n_{j}}{x}\binom{N-n_{j}}{k-x}}{\binom{N}{k}}
$$

## Sequential, Uniform, Distinct

## Sequential, Uniform, Distinct

- Estimate the distribution of distance between two qualifying tuples
- Bitvector $B, b$ bits are set to 1
- First, let's find the distribution of number of zeros
- before first 1
- between two consecutive 1s
- after last 1
- $B-j-1$ positions for $i$
- every bitvector has $b-1$ sequences of a form $10 \ldots 01$
$-\frac{(B-j-1)\binom{B-j-2}{b-2}}{(b-1)\binom{B}{b}}=\frac{\binom{B-j-1}{b-1}}{\binom{B}{b}}$
- now, the expected number of $0 \mathrm{~s}: \frac{B-b}{b+1}$
- then, the expected total number of bits between first and last 1 :


## Sequential, Uniform, Distinct

- Estimate the distribution of distance between two qualifying tuples
- Bitvector $B, b$ bits are set to 1
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- $B-j-1$ positions for $i$
- every bitvector has $b-1$ sequences of a form $10 \ldots 01$
$-\frac{(B-j-1)\binom{B-j-2}{b-2}}{(b-1)\binom{B}{b}}=\frac{\binom{B-j-1}{b-1}}{\binom{B}{b}}$
- now, the expected number of $0 \mathrm{~s}: \frac{B-b}{b+1}$
- then, the expected total number of bits between first and last 1: $B-\frac{B-b}{b+1}=\frac{B b+b}{b+1}$


## Histograms

A histogram $H_{A}: B \rightarrow \mathbb{N}$ over a relation $R$ partitions the domain of the aggregated attribute $A$ into disjoint buckets $B$, such that

$$
\begin{aligned}
& \qquad H_{A}(b)=|\{r \mid r \in R \wedge R . A \in b\}| \\
& \text { and thus } \sum_{b \in B} H_{A}(b)=|R| .
\end{aligned}
$$

## Histograms

A rough histogram might look like this:


## Using Histograms (3)

Given a histogram, we can approximate the selectivities as follows:

$$
\begin{array}{ll}
A=c & \frac{\sum_{b \in B: c \in b} H_{A}(b)}{\sum_{b \in B} H_{A}(b)} \\
A>c & \frac{\sum_{b \in B: c \in b} \frac{\max (b)-c}{\max (b)-\min (b)} H_{A}(b)+\sum_{b \in B: \min (b)>c} H_{A}(b)}{\sum_{b \in B} H_{A}(b)} \\
A_{1}=A_{2} & \frac{\sum_{b_{1} \in B_{1}, b_{2} \in B_{2}, b^{\prime}=b_{1} \cap b_{2}: b^{\prime} \neq \emptyset} \frac{\max \left(b^{\prime}\right)-\min \left(b^{\prime}\right)}{\max \left(b_{1}\right)-\min \left(b_{1}\right)} H_{A_{1}}\left(b_{1}\right) \frac{\max \left(b^{\prime}\right)-\min \left(b^{\prime}\right)}{\max \left(b_{2}\right)-\min \left(b_{2}\right)} H_{A_{2}}\left(b_{2}\right.}{\sum_{b_{1} \in B_{1}} A_{A_{1}}\left(b_{1}\right) \sum_{b_{2} \in B_{2}} H_{A_{2}}\left(b_{2}\right)}
\end{array}
$$

## Exam: Algorithms

- Exact vs approximate
- Deterministic vs probabilistic

Other important divisions:

- Bottom-up vs top-down (DP vs memoization)
- Random vs pseudo-random tree generation
- Hybrid algos

Some heuristics can be combined (and some can't):

- GOO + Iterative improvement
- Iterative Improvement + Simulated Annealing
- Does it make sense to do SA and then II?


## Exam: Algorithms

Important aspects:

- When can it be applied?
- When is it good?
-What's the runtime complexity?


## Exam: Formulas

Important ones include, but are not limited to:

- Cost functions!
- rank in IKKBZ
- benefit in query simplification
- Yao formula (unless you can derive it yourself quick)
- Histograms
- ...
- Exam: Hörsaal 2
- 30 July, 08:30-10:00
- No repeat exam

