Low Latency Query Planning and Processing

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Hardware gets fast

- Large main memory
- Fast SSDs
- Many core machines

Low latency queries

- Still bound by CPU capabilities
- Algorithmic changes
Algorithmic challenges

Query processing
  ● Intra-query parallelism
  ● Shared state

Query planning
  ● Cardinality estimation
  ● Algebra optimization
Groupjoin – Idea

- Combined execution of compatible join and aggregation
- Q: “Total sales per customer”

```
join = Hashtable()
for c in customer:
    join[c.id] = c

group = Hashtable()
for o in orders:
    if join.contains(o.c_id):
        group[o.c_id].sum += o.price

groupjoin = Hashtable()
for c in customer:
    groupjoin[c.id] = c

for o in orders:
    if groupjoin.contains(o.c_id):
        groupjoin[o.c_id].sum += o.price
```
Groupjoin – Avoiding contention

- Shared hash table unsuitable for multithreaded execution
- Four execution strategies for parallel groupjoin:

  - Separate
  - Eager
  - Memoizing
  - Index
Nested Aggregates

- Common in analytical queries
- HAVING predicates are hard to estimate

Q: “Large orders”

```sql
select l_orderkey
from lineitem
group by l_orderkey
having sum(l_quantity) > 300
```
Nested Aggregates

- Common in analytical queries
- HAVING predicates are hard to estimate
- But have significant impact on execution plans
Estimating Aggregates

- Numerical columns $\sim \mathcal{N}(\mu, \sigma^2)$
- Cheap and generalizes nicely, but inherently symmetric
Estimating Aggregates

- Using a skew-normal distribution
Practical Groupjoins and Nested Aggregates

- Effects ⅛ of queries
- +23% in TPC-H, +6% in TPC-DS
Query Optimization with Indexed Algebra

- Complex queries on small workloads
  - BigQuery: 90% of queries processed less than 100 MB of data
  - Tableau Public: 90% of workbooks are less than 100k tuples

- TPC-H
  - Scale 1: 0.8 ms optimization, 20 ms execution
  - Scale 0.01: 0.8 ms optimization, 0.2 ms execution

- Optimization time scales super-linear with query complexity
Algebra

● Relational algebra trees
  ○ Operators
  ○ Expressions
  ○ Columns / IUs

● Analyze data-flow for optimization
  ○ Which path?
  ○ Modifications?
  ○ Materialized?
Optimization

- Reason about the algebra to derive optimization possibilities
- Top-down, operator at a time
  - Needs $O(n^2)$ column sets
- Path-centric
  - Still $O(n)$ length
  - With indexing: $O(\log n)$
Indexing Algebra

- Index paths through the algebra
  - Faster path traversal
- Binary search trees on path depth
- Paths from root overlap
- Link/cut trees support that efficiently
Indexing Algebra

- Index paths through the algebra
  ➡ Faster path traversal
- Binary search trees on path depth
- Paths from root might overlap
- **Link/cut trees** support that efficiently

<table>
<thead>
<tr>
<th>Rel. Algebra</th>
<th>Transformation</th>
<th>Traversal</th>
</tr>
</thead>
<tbody>
<tr>
<td>w/o index</td>
<td>$O(1)$</td>
<td>$O(n)$</td>
</tr>
<tr>
<td>static index</td>
<td>$O(n)$</td>
<td>$O(\log n)$</td>
</tr>
<tr>
<td>path labeling</td>
<td>$O(n)$</td>
<td>$O(1)$</td>
</tr>
<tr>
<td>Indexed Algebra</td>
<td>$O(\log n)$</td>
<td>$O(\log n)$</td>
</tr>
</tbody>
</table>
Indexed Algebra Performance

- Significant overall improvements
- 10 - 30% faster optimization
- 8% better end-to-end latency in Tableau Public
Conclusion

Query processing
- ✅ Intra-query parallelism
- ✅ Shared state

Query planning
- ✅ Cardinality estimation
- ✅ Algebra operations