

#### Problem:

- want to sort an arbitrary amount of data stored on disk
- accessing disk is slow so we do not want to access each value individually
- sorting in main memory is fast but main memory size is limited

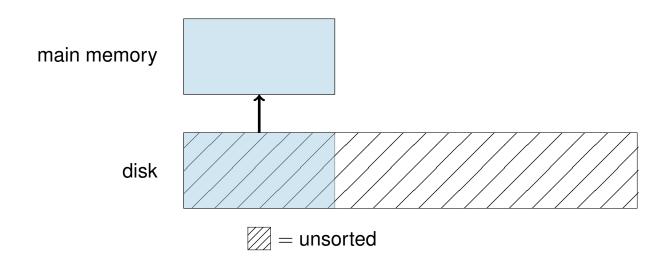
#### Solution:

- load pieces (called "runs") of the data into main memory and sort them
- with m values fitting into main memory and d values that should be sorted this results in  $k = \lceil \frac{d}{m} \rceil$  sorted runs
- do a k-way merge of all runs

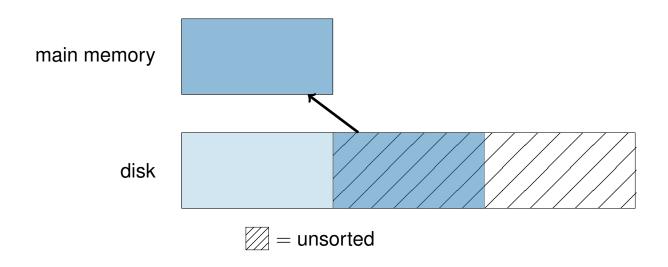


main memory	
disk	
	= unsorted

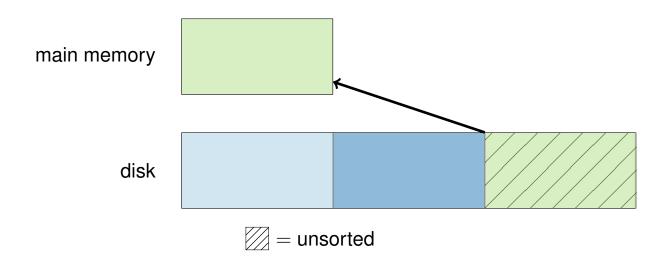




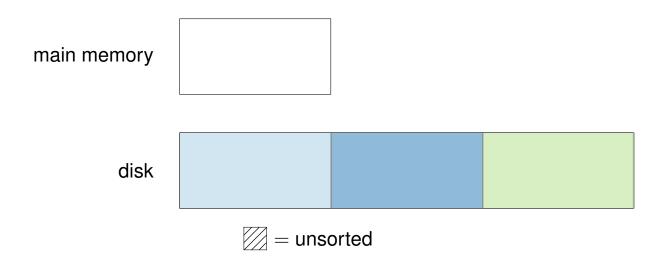




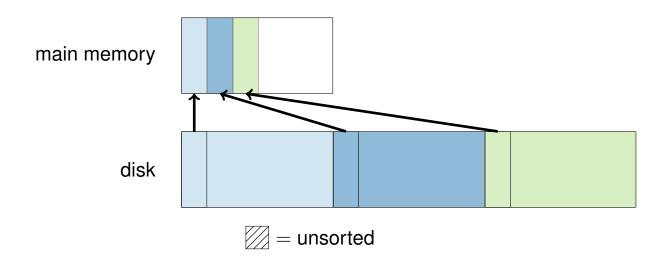




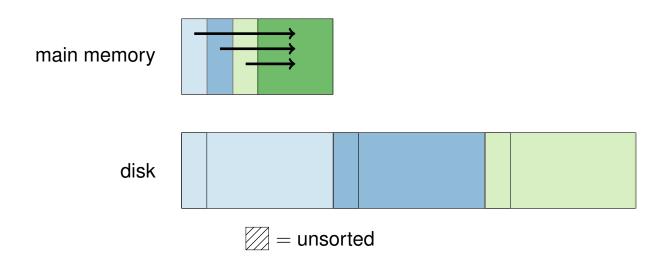




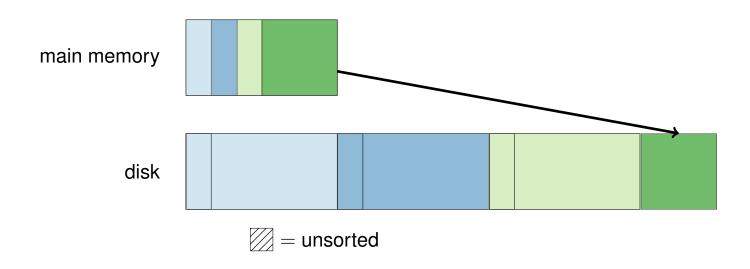














main memory			
disk			
	= unso	orted	

Step: done