

# Query Optimization

## Exercise Session 11

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February 6, 2017

## Combinatorics 101

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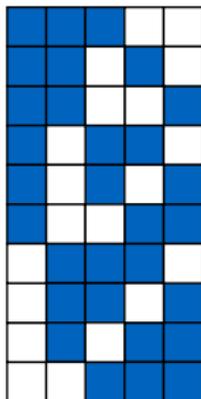
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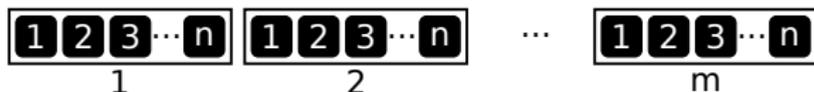
Example: Choose 3 out of 5:  $\binom{5}{3} = \frac{5!}{2! \cdot 3!} = \frac{120}{2 \cdot 6} = 10$



Direct, Uniform, Distinct

## Waters/Yao Bottom-Up

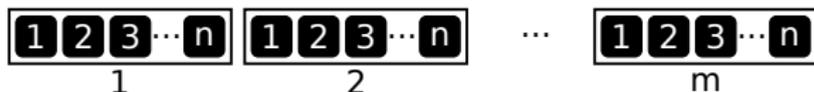
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- ▶ How many distinct subsets of size  $k$  exist?

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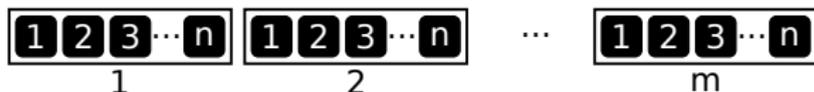
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- ▶ How many distinct subsets of size  $k$  exist?  $\binom{N}{k}$
- ▶ How many distinct subsets of size  $k$  exist, where a page does not contain any chosen tuples? Choose  $k$  from all but one page, i.e. from  $N - n$  tuples:

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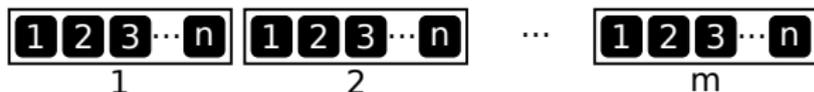
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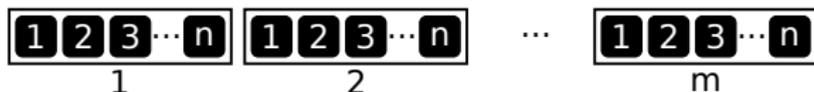
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$$p := \frac{\binom{N-n}{k}}{\binom{N}{k}}$$

- ▶ What is the probability that a certain page contains at least one tuple?  $1 - p$ ... unless all pages have to be involved ( $k > N - n$ ).
- ▶ Multiplied by the number of pages, we get the number of qualifying pages, denoted  $\bar{y}_n^{N,m}(k)$ .

# Approximation

Let  $m = 50$ ,  $n = 1000 \Rightarrow N = 50k$ ,  $k = 100$

$$\text{Yao (exact)} : p = \frac{\binom{N-n}{k}}{\binom{N}{k}} = \prod_{i=0}^{k-1} \frac{N-n-i}{N-i} = \prod_{i=0}^{99} \frac{49k-i}{50k-i} = 13.2\%$$

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$$\text{Waters} : p \approx \left(1 - \frac{k}{N}\right)^n \approx 13.5\%$$

Direct, Uniform, Non-Distinct

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 $f : (x_1, x_2, \dots, x_k) \mapsto (x_1 + 0, x_2 + 1, \dots, x_k + (k - 1))$

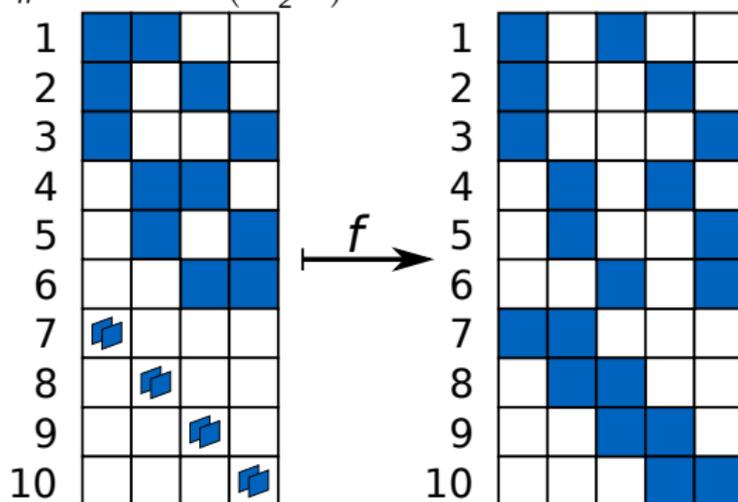
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- ▶ Example: Choose 2 from 4

- ▶ # sets:  $\binom{4}{2}$
- ▶ # multisets:  $\binom{4+2-1}{2}$



- ▶ Like Yao, but not necessarily distinct
- ▶ Same formula as Yao, but:
  - ▶ We don't need to distinguish cases when computing the probability that a bucket contains at least one item
  - ▶ We substitute  $N$  by  $N + k - 1$  to compute  $\tilde{p}$

Sequential, Uniform, Distinct

## Sequential, Uniform, Distinct

- ▶ Estimate the distribution of distance between two qualifying tuples
- ▶ Bitvector  $B$ ,  $b$  bits are set to 1
- ▶ First, let's find the distribution of number of zeros
  - ▶ before first 1
  - ▶ between two consecutive 1s
  - ▶ after last 1
- ▶  $B - j - 1$  positions for  $i$
- ▶ every bitvector has  $b - 1$  sequences of a form  $10 \dots 01$
- ▶ 
$$\frac{(B-j-1)\binom{B-j-2}{b-2}}{(b-1)\binom{B}{b}} = \frac{\binom{B-j-1}{b-1}}{\binom{B}{b}}$$
- ▶ now, the expected number of 0s:  $\frac{B-b}{b+1}$
- ▶ then, the expected total number of bits between first and last 1:

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- ▶ then, the expected total number of bits between first and last 1:  $B - \frac{B-b}{b+1} = \frac{Bb+b}{b+1}$

# Histograms

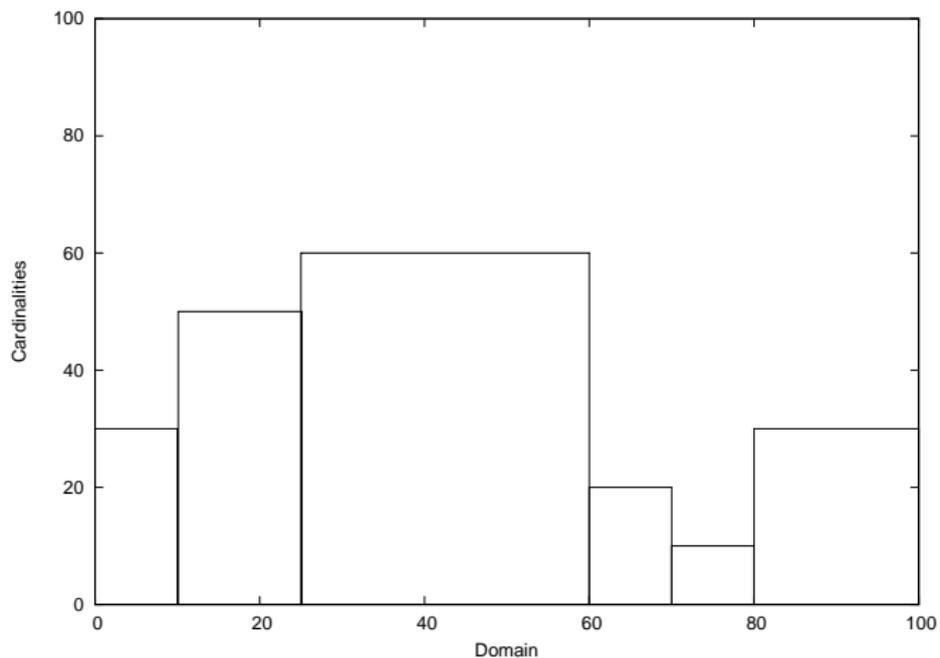
A histogram  $H_A : B \rightarrow \mathbb{N}$  over a relation  $R$  partitions the domain of the aggregated attribute  $A$  into disjoint buckets  $B$ , such that

$$H_A(b) = |\{r \mid r \in R \wedge R.A \in b\}|$$

and thus  $\sum_{b \in B} H_A(b) = |R|$ .

# Histograms

A rough histogram might look like this:



## Using Histograms (3)

Given a histogram, we can approximate the selectivities as follows:

$$A = c \quad \frac{\sum_{b \in B: c \in b} H_A(b)}{\sum_{b \in B} H_A(b)}$$

$$A > c \quad \frac{\sum_{b \in B: c \in b} \frac{\max(b) - c}{\max(b) - \min(b)} H_A(b) + \sum_{b \in B: \min(b) > c} H_A(b)}{\sum_{b \in B} H_A(b)}$$

$$A_1 = A_2 \quad \frac{\sum_{b_1 \in B_1, b_2 \in B_2, b' = b_1 \cap b_2: b' \neq \emptyset} \frac{\max(b') - \min(b')}{\max(b_1) - \min(b_1)} H_{A_1}(b_1) \frac{\max(b') - \min(b')}{\max(b_2) - \min(b_2)} H_{A_2}(b_2)}{\sum_{b_1 \in B_1} H_{A_1}(b_1) \sum_{b_2 \in B_2} H_{A_2}(b_2)}$$

## Exercise

Given a relation with 3 pages and two tuples per page, compute the average number of accessed pages when reading 2 tuples.

## Exercise

Given the following Histogram for  $R.A$ :

bucket	[0, 20)	[20, 40)	[40, 60)	[60, 80)	[80, 100)
count	2	4	0	2	2

Estimate how many tuples the following SQL query will fetch

```
select *  
from R  
where R.A < 30
```

## Exam: Algorithms

- ▶ exact vs. approximate
- ▶ deterministic vs. probabilistic

Other important divisions:

- ▶ bottom-up vs. top-down (DP vs. memoization)
- ▶ random vs. pseudo-random trees
- ▶ meta-heuristics and hybrid algorithms

Some heuristics can be combined

- ▶ GOO + Iterative Improvement
- ▶ Iterative Improvement + Simulated Annealing
- ▶ Does it make sense to do SA and then II?

# Exam: Algorithms

Important aspects:

- ▶ When can it be applied?
- ▶ When is it good?
- ▶ What is the runtime complexity?

## Exam: Formulas

Important ones include, but are not limited to:

- ▶ cost functions
- ▶ rank in IKKBZ
- ▶ ordering benefit in query simplification
- ▶ Yao, etc.
- ▶ Histograms

# Exam

- ▶ Hörsaal 1 (00.02.001)
- ▶ 27.02.2017 13:30 - 15:00
- ▶ no repeat exam

# Bonus

In no particular order:

- ▶ Matthias Adams and David Werner
- ▶ Maximilian Bandle and Alexander Beischl
- ▶ Kordian Bruck and Philipp Fent
- ▶ Florian Dreier and Shwetha Suresh
- ▶ Dominik Durner and Moritz Sichert
- ▶ Frank Hermann and Phuoc Le
- ▶ Hans Kirchner and Fabian Schurig
- ▶ Julius Michaelis and Maxi Weininger
- ▶ Michael Schreier and Sebastian Stein